The Statistical Mechanics of Formal Privacy

Theory and Experiments

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The views expressed herein are those of the author and not necessarily those of the U.S. Census Bureau



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Overview:

- The need for Disclosure Avoidance and how it's done
- The perspective offered by Statistical Mechanics in the context of privacy
- The relevant elements of Statistical Mechanics.
 - Simulated Annealing (SA) and Markov Chain Monte Carlo (MCMC) simulations
 - Changing perspective: DP noise injection vs. DP noise evolution.
 - Unsolving optimization problems as a noise-evolution method (as opposed to a 'noise-injection' method).
- The Boltzmann Machine Privacy Framework (BMPF)
 - Description of consensus functions for histograms
 - Generating 'candidate' histograms
 - How the BMPF satisfies ($\epsilon(t)$, $\delta(k)$)-DP
- Experimental Results
 - MCMC burn-in issues, ways to ameliorate this *i.e.*, stopping criteria.
- Conclusion



Disclosure Avoidance and What it Entails

- Can't publish actual datasets collected by CB
 - Privacy laws prohibit disclosure. E.g., Title 13, Title 26 and others.
 - Publishing data with many traditional DA methods can still disclose private information (e.g., Governor Weld's medical records were 'anonymized')
 - But good public policy-making requires *some* use of the data.
- How can we publish data yet maintain privacy?
- Differential Privacy (2006), formerly referred to as *epsilon indistinguishability* provides a methodology that guarantees a quantifiable level of privacy via a 'privacy budget'.
- It infuses data with 'calibrated noise' to achieve this quantifiable level of privacy.



Examples:

- Histograms are a common type of dataset developed at Census
- They reflect counts of entities (people) that are associated with certain mutually exclusive combinations of attributes
- Publishing actual counts can lead to complete privacy loss
- DP modifies these counts in a probabilistic manner such that there is a quantifiable level of privacy yet still maintains usability/utility.



Definition of Differential Privacy

$\Pr\{\mathcal{M}(\mathbf{x}')\} \in S\} \le e^{\epsilon} \Pr\{\mathcal{M}(\mathbf{x}) \in S\} + \delta$

E.g., adding random noise to pixels in a picture to blur the faces of people in the picture making it hard to identify the people in the picture, yet enabling a fairly accurate counting of the number of people.

Quantifies the fundamental tradeoff between accuracy and privacy.

Lots of way to create 'noisy data':

- 1. Add random variates to the actual data.
 - 2. Consider the actual data as the 'optimum data' in an optimization problem and produce *sub-optimum* data.



Simulated Annealing and Markov Chains

Simulated Annealing (SA) circa 1983 is a meta-heuristic that can 'solve' a wide variety of optimization problems.

Hallmarks:

• Based on the Metropolis Algorithm (an accept/reject method), it enables Markov Chain Monte Carlo (MCMC) sampling. The MAC: Let $\Delta E = E_{cand} - E_{curr}$ $Pr\{Accept E_{cand}\} = \begin{cases} e^{-\Delta E/t} \text{ if } \Delta E \ge 0\\ 1 & \text{otherwise} \end{cases}$ $Pr\{Accept Candidate j\} = \frac{1}{1 + e^{-\Delta f_{ji}/t}}$ $\pi_i p_{ij} = \pi_j p_{ji}$

$$\pi_i(t) = \frac{e^{-E_i/t}}{\sum_j e^{-E_j/t}}$$

• SA converges in probability under WLLN to the *globally optimal* solution:



$$\lim_{t \to 0^+} \pi_i(t) = \frac{1}{|\mathcal{C}_{\texttt{opt}}|}$$

Simulated Annealing

- Used to solve a wide variety of optimization problem by virtue of its simplicity, convergence properties and generalizability.
- Requirements:
 - 1. A well defined configuration space.
 - 2. A well defined objective function.
 - 3. A 'good' candidate generation scheme.
 - 4. An appropriate 'cooling schedule'. (We'll just need a fixed temperature.)
- MA/SA effectively moves (transitions) from one configuration to another under the influence of these four elements.
- Transition probabilities

🕨 Markov Chain

Instead of using SA to find the global optimum (the 'true' configuration), we use it to *move away* from the optimal solutions to find a *sub-optimal* configuration which is equivalent to a 'noisy' configuration by holding the temperature to some positive value.



Noise-Injection Paradigms



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Stationary Distributions

$$\mathcal{C} = \{(2,0,0), (1,1,0), (1,0,1), (0,2,0), (0,1,1), (0,0,2)\}$$



Distributions of the Configuration Space (C) of all histograms with k = 3 bins and N = 2 individuals

For a fixed N and fixed k, the size of the configuration space = $\binom{N+k}{k-1}$

$$\binom{N+k-1}{k-1} = \binom{4}{2} = 6$$



E.g., if N = 50, $k = 5 \rightarrow |\mathcal{C}| = 316251$

The Main Idea Behind Data Noise Injection Using The Metropolis/Simulated Annealing Approach

- 'Ground truth' = 'observed data' = 'sampled data' = 'optimal data'
- 'Disclosed data' = 'noise injected data' = 'suboptimal data'

Ground truth/optimal histogram:[1,2,3,20,24]Total count = 50t = 10t = 100[1,4,5,19,21][10,8,5,11,16][3,5,6,16,24][6,7,10,13,14][3,4,4,18,21][7,10,11,9,13]



Notice that the total count in all these histograms = 50

Candidate Generation Mechanism



Histogram

- Negative bin values are impossible.
- Invariants correspond to bins that are ignored.
- Total counts are unchanged.



Can the Metropolis/SA Algorithm Create Markov Chains that are DP?

Define the *configuration space*:

Definition : Define $\mathbf{x} \in \mathbb{N}^n$ and corresponds to a histogram vector of n counts.

Definition : Given a vector $\mathbf{x} \in \mathbb{N}^n$ and where $\sum_{i=1}^n x_i = N$, define

$$\mathcal{C}_{\mathbf{x}} = \{ \mathbf{x} \in \mathbb{N}^n : \forall \mathbf{x} \in \mathbb{N}^n, \sum_{i=1}^N x_i = N \}$$



Can the Metropolis/SA Algorithm Create Markov Chains that are DP?

Define the following objective functions for neighboring configurations **x** and **y**:

$$c_{\mathbf{x}}(\mathbf{z}) = -(\mathbf{z} - \mathbf{x})^{\mathsf{T}} \mathbf{W}_{\mathbf{x}}(\mathbf{z} - \mathbf{x})$$
 and

Definition : For all neighboring configurations $\mathbf{x}, \mathbf{y} \in C$, define the global sensitivity

$$s = \max_{\mathbf{z} \in \mathcal{C}} \max_{\mathbf{x}, \mathbf{y}: d(\mathbf{x}, \mathbf{y}) \le 2} \|c_{\mathbf{x}}(\mathbf{z}) - c_{\mathbf{y}}(\mathbf{z})\|_{2}$$

Each dataset/configuration **x** and **y** induces its own Markov Chain:





The Boltzmann Machine Mechanism

Definition : Define the Boltzmann Machine Mechanism $\mathcal{B}_{k,t}(\mathbf{x})$ as the random configuration of a discrete time, irreducible and aperiodic Markov chain generated by application of the Metropolis algorithm at temperature t after k iterations given a dataset \mathbf{x} where \mathbf{x} is the initial configuration of the Markov Chain. This simply corresponds to the k^{th} -step transition probability of the Markov Chain. Thus,

$$\Pr\{\mathcal{B}_{k,t}(\mathbf{x}) = \mathbf{z}\} \equiv \Pr\{X_k(t) = \mathbf{z} | X_0(t) = \mathbf{x}\} \equiv p_{\mathbf{x},\mathbf{z}}^{(k)}$$

Theorem : Let $(X_k(t))_{k\geq 0}$ be an irreducible, aperiodic Markov Chain based on the Metropolis algorithm as in Theorem 4. Then the BMM $\mathcal{B}_{k,t}(\mathbf{x})$ satisfies $(\epsilon(t), \delta(k)) - DP$ where $\delta(k) \to 0$ as $k \to \infty$.



Markov Chain Convergence

Theorem : Let X_k be an irreducible and aperiodic Markov Chain on a finite configuration space C with stationary distribution $\pi(t)$ at fixed temperature t. Then there exist constants $\alpha \in (0, 1)$ and C > 0 such that for all state vectors $\mathbf{v}^{[k]}$

$$\max_{\mathbf{v}^{[0]} \in \mathcal{C}} \|\mathbf{v}^{[k]} - \pi(t)\|_{TV} \le C\alpha^k \tag{10}$$

 $\Pr\{\mathcal{B}_{k,t}(\mathbf{x}) \in \mathcal{S}\} \le e^{\epsilon(t)} \Pr\{\mathcal{B}_{k,t}(\mathbf{y}) \in \mathcal{S}\} + \delta(k)$

where
$$\epsilon(t) = \frac{2s}{t}$$
 and $\delta(k) \to 0$ as $k \to \infty$

As $t \to \infty$, $\epsilon(t) \to 0$ As $t \to 0$, $\epsilon(t) \to \infty$



Convergence of Two or More Markov Chains





Convergence of Two or More Markov Chains

Theorem:

Given p irreducible, aperiodic Markov Chains $X_m^{[k]}, m = 1, 2, ..., p$ and a metric space $d(\cdot, \cdot)$ where

$$M_m^{[k]} = d\left(f(X_m^{[k]}), f(\mathbf{x}_{\text{OPT}})\right)$$

where $\forall m, M_m^{[k]} \xrightarrow{\text{Prob}} 0$ as $k \to \infty$ and some function

$$Y^{[k]} = g(X_1^{[k]}, X_2^{[k]}, \dots, X_p^{[k]})$$

such that

 $Y^{[k]} \to 0$ as $k \to \infty$

and

$$Y^{[k]} = 0 \iff f(X_1) = f(X_2) = \ldots = f(X_p)$$

then for any chain m and k sufficiently high

$$\Pr\{M_m^{[k]} = 0 | Y^{[k]} = 0\} > \Pr\{M_m^{[k]} = 0\}.$$



Convergence of Two or More Markov Chains

Corollary: Let p be the number of independent Markov Chains and let $Y^{[k,p]}$ be the metric among the p chains as defined above at time index k. Then for any chain m with $p \ge 2$ and k sufficiently large, then

$$\Pr\{M_m^{[k]} = 0 | Y^{[k,p+1]} = 0\} > \Pr\{M_m^{[k]} = 0 | Y^{[k,p]} = 0\} > \Pr\{M_m^{[k]} = 0\}.$$



Stopping Criteria:

How can we apply the foregoing theorems when we are not converging to the 'true' data?

We can define a random variable that converges to 0 based on the ergodic theorem:

$$\frac{1}{n} \sum_{k=0}^{n-1} c_{\mathbf{x}}(X_{1}^{[k]}) \xrightarrow{a.s.} \sum_{\mathbf{z} \in \mathcal{C}} \pi_{\mathbf{z}|\mathbf{x}}(t) c_{\mathbf{x}}(\mathbf{z}) \text{ as } n \to \infty \qquad \text{Convergence to the expected objective function value.}$$

$$\lim_{k \to \infty} \left| \frac{1}{n} \sum_{n=0}^{n-1} c_{\mathbf{x}}(X_{1}^{[k]}) - \sum_{\mathbf{z} \in \mathcal{C}} \pi_{\mathbf{z}|\mathbf{x}}(t) c_{\mathbf{x}}(\mathbf{z}) \right| = 0$$

$$M_{1}^{[k]} \equiv \frac{1}{n} \sum_{n=0}^{n-1} c_{\mathbf{x}}(X_{1}^{[k]}) - \sum_{\mathbf{z} \in \mathcal{C}} \pi_{\mathbf{z}|\mathbf{x}}(t) c_{\mathbf{x}}(\mathbf{z}) \to 0$$

$$M_{2}^{[k]} \equiv \frac{1}{n} \sum_{n=0}^{n-1} c_{\mathbf{x}}(X_{2}^{[k]}) - \sum_{\mathbf{z} \in \mathcal{C}} \pi_{\mathbf{z}|\mathbf{x}}(t) c_{\mathbf{x}}(\mathbf{z}) \to 0$$

$$Y^{[k]} = M_{1}^{[k]} - M_{2}^{[k]} \xrightarrow{\text{prob}} 0$$

Experiments: *t* = 10, 50; Run Length = 5000; Replications = 30





Histogram: (1, 2, 3, 20, 24)

Conclusion

- Metropolis/Simulated Annealing Algorithm can serve as a basis for noise injection.
- Boltzmann Machine Mechanism can provide 'noise-evolution'.
- Constructing candidate generation mechanisms can obviate the need for post-processing.
- Resulting Markov Chains satisfies $(\epsilon(t), \delta(k)) DP$.
- Flexibility in implementation: data object types and objective functions.
- Possibility of *tuning* sensitivity by modifying the weight matrices.
- Bears some striking similarities to the Exponential Mechanism.



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